

A Novice-Friendly Induction Tactic for Lean

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```
case semantics.big_step.skip
t h : state
⊢ false
```

```
case semantics.big_step.assign
t : state,
h_x : string,
h_a : state → ℕ,
h_s : state
⊢ false
```

```
case semantics.big_step.seq
t : state,
h_S h_T : program,
h_s h_t h_u : state,
h_hS : (h_S, h_s) ⇒ h_t,
h_hT : (h_T, h_t) ⇒ h_u,
h_ih_hS h_ih_hT : false
⊢ false
```

```
case semantics.big_step.ite_true
t : state,
h_b : state → Prop,
h_S h_T : program,
h_s h_t : state,
h_hcond : h_b h_s,
h_hbody : (h_S, h_s) ⇒ h_t,
h_ih : false
⊢ false
```

```
case semantics.big_step.ite_false
t : state,
h_b : state → Prop,
h_S h_T : program,
h_s h_t : state,
h_hcond : ¬h_b h_s,
h_hbody : (h_T, h_s) ⇒ h_t,
h_ih : false
⊢ false
```

```
case semantics.big_step.while_true
t : state,
h_b : state → Prop,
h_S : program,
h_s h_t h_u : state,
h_hcond : h_b h_s,
h_hbody : (h_S, h_s) ⇒ h_t,
h_hrest : (while h_b h_S, h_t) ⇒ h_u,
h_ih_hbody h_ih_hrest : false
⊢ false
```

```
case semantics.big_step.while_false
t : state,
h_b : state → Prop,
h_S : program,
h_s : state,
h_hcond : ¬h_b h_s
⊢ false
```

Complex Indices

Induction Hypothesis Generalisation

Naming

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Complex Indices

$\forall S s t,$
 $(\text{while } (\lambda _ , \text{true}) S, s) \Rightarrow t \rightarrow \text{false}$

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$\forall S s t,$
 $(\text{while } (\lambda _, \text{true}) S, s) \Rightarrow t \rightarrow \text{false}$

case skip
 $t t' : \text{state}$
 $\vdash \text{false}$

McBride's Method

...
h : (while (λ _, true) S, s) \Rightarrow t
 \vdash false

McBride's Method

...
h : (while (λ _, true) S, s) \Rightarrow t
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...
i : program \times state,
h : i \Rightarrow t
 \vdash (while (λ _, true) S, s) = i \rightarrow false

McBride's Method

...
 $h : (\text{while } (\lambda _ , \text{true}) S, s) \Rightarrow t$
 $\vdash \text{false}$

...
 $i : \text{program} \times \text{state},$
 $h : i \Rightarrow t$
 $\vdash (\text{while } (\lambda _ , \text{true}) S, s) = i \rightarrow \text{false}$

case skip

...
 $\vdash (\text{while } (\lambda _ , \text{true}) S, s) = (\text{skip}, u) \rightarrow \text{false}$

Induction Hypothesis Simplification

$\forall S' s',$
 $(\text{while } (\lambda _ , \text{true}) S', s') =$
 $(\text{while } (\lambda _ , \text{true}) S, t) \rightarrow$
 false

Induction Hypothesis Simplification

$\forall S' s',$
 $(\text{while } (\lambda _ , \text{true}) S', s') =$
 $(\text{while } (\lambda _ , \text{true}) S, t) \rightarrow$
 false

$(\text{while } (\lambda _ , \text{true}) ?S', ?s') =$
 $(\text{while } (\lambda _ , \text{true}) S, t) \rightarrow$
 false

Induction Hypothesis Simplification

$\forall S' s',$
 $(\text{while } (\lambda _ , \text{true}) S', s') =$
 $(\text{while } (\lambda _ , \text{true}) S, t) \rightarrow$
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$(\text{while } (\lambda _ , \text{true}) ?S', ?s') =$
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Induction Hypothesis Simplification

$\forall S' s',$
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$(\text{while } (\lambda _ , \text{true}) ?S', ?s') =$
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 false

false

Complex Indices

Induction Hypothesis Generalisation

Naming

Induction Hypothesis Generalisation

[Benjamin C. Pierce et al., Software Foundations]

Induction Hypothesis Generalisation

$$\forall n\ m, n + n = m + m \rightarrow n = m$$

[Benjamin C. Pierce et al., Software Foundations]

Induction Hypothesis Generalisation

$$\forall n\ m, n + n = m + m \rightarrow n = m$$

case succ

$m\ n : \mathbb{N}$

$ih : n + n = m + m \rightarrow n = m$

$h : n + n + 2 = m + m$

$\vdash n + 1 = m$

Induction Hypothesis Generalisation

$$\forall n m, n + n = m + m \rightarrow n = m$$

case succ

$$m n : \mathbb{N}$$

$$ih : n + n = m + m \rightarrow n = m$$

$$h : n + n + 2 = m + m$$

$$\vdash n + 1 = m$$

...

$$ih : \forall m, n + n = m + m \rightarrow n = m$$

$$h : n + n + 2 = m + m$$

$$\vdash n + 1 = m$$

$x : X$

$n \ m : \mathbb{N}$

$\vdash n + m = m + n$

$x : X$

$n\ m : \mathbb{N}$

$\vdash n + m = m + n$

$ih : X \rightarrow \forall m, n + m = m + n$

Criteria for Auto-Generalisation

Connection with the target

$n\ m : \mathbb{N}$

$\vdash n + m = m + n$

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Connection with the target

$n \ m : \mathbb{N}$

$\vdash n + m = m + n$

Connection with the eliminated hypothesis

$n : \mathbb{N}$

$h_1 : n \geq 0$

$h_2 : n \neq 0$

$\vdash n > 0$

Complex Indices

Induction Hypothesis Generalisation

Naming

Naming

Naming

$\forall \alpha (r : \alpha \rightarrow a \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

Naming

$\forall \alpha (r : \alpha \rightarrow a \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

$\alpha : \text{Type}$

$r : \alpha \rightarrow \alpha \rightarrow \text{Type}$

$c\ a\ b\ h_{1_x}\ h_{1_y}\ h_{1_z} : \alpha$

$h_{1_hr} : r\ h_{1_x}\ h_{1_y}$

$h_{1_ht} : \text{tc } r\ h_{1_y}\ h_{1_z}$

$h_{1_ih} : \text{tc } r\ h_{1_z}\ c \rightarrow \text{tc } r\ h_{1_y}\ c$

$h_2 : \text{tc } r\ h_{1_z}\ c$

$\vdash \text{tc } r\ h_{1_x}\ c$

Naming

$\forall \alpha (r : \alpha \rightarrow \alpha \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

$\alpha : \text{Type}$
 $r : \alpha \rightarrow \alpha \rightarrow \text{Type}$
 $c\ x\ y\ z : \alpha$
 $hr : r\ x\ y$
 $h_1 : \text{tc } r\ y\ z$
 $IHh_1 : \text{tc } r\ z\ c \rightarrow \text{tc } r\ y\ c$
 $h_2 : \text{tc } r\ z\ c$
 $\vdash \text{tc } r\ x\ c$

Naming

$\forall \alpha (r : \alpha \rightarrow \alpha \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

$\alpha : \text{Type}$
 $r : \alpha \rightarrow \alpha \rightarrow \text{Type}$
 $a\ y\ b\ c : \alpha$
 $hr : r\ a\ y$
 $h_1 : \text{tc } r\ y\ b$
 $ih : \forall c, \text{tc } r\ b\ c \rightarrow \text{tc } r\ y\ c$
 $h_2 : \text{tc } r\ b\ c$
 $\vdash \text{tc } r\ a\ c$

Index Association

h_1 : t c r a b

...

hr : r x y

h_1 : t c r y z

...

Index Association

$h_1 : tc\ r\ a\ b$

...

$hr : r\ x\ y$

$h_1 : tc\ r\ y\ z$

...

```
inductive tc { $\alpha$  : Type} (r :  $\alpha \rightarrow \alpha \rightarrow$  Type) :  
   $\alpha \rightarrow \alpha \rightarrow$  Type  
| base :  $\forall$  x y (hr : r x y), tc x y  
| step :  $\forall$  x y z (hr : r x y) (ht : tc y z),  
  tc x z
```

Type-Based Naming

```
class variable_names ( $\alpha$  : Type) : Type :=  
(names : list name)
```

Type-Based Naming

```
class variable_names ( $\alpha$  : Type) : Type :=  
(names : list name)
```

```
instance : variable_names  $\mathbb{N}$  :=  
<['n, 'm, 'o]>
```

```
instance : variable_names Prop :=  
<['P, 'Q, 'R]>
```

```
instance { $\alpha$ } [variable_names  $\alpha$ ] :  
  variable_names (list  $\alpha$ ) :=  
<(names  $\alpha$ ).map ( $\lambda$  n, n.append_suffix "s")>
```


Summary

- ▶ Simplification of induction hypotheses
- ▶ Auto-generalisation
- ▶ Index- and type-based naming
- ▶ (Not in this talk:) Evaluation of Lean 3's Metaprogramming Framework

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