

A Novice-Friendly Induction Tactic for Lean (Short Talk)

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Simplification of Induction Hypotheses

[Conor McBride, Elimination with a Motive, TYPES 2000]

Simplification of Induction Hypotheses

$\forall S' s',$
 $(\text{while } (\lambda _ , \text{true}) S', s') =$
 $(\text{while } (\lambda _ , \text{true}) S, t) \rightarrow$
 false

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Automatic Generalisation of Induction Hypotheses

$$\forall n\ m, n + n = m + m \rightarrow n = m$$

[Benjamin C. Pierce et al., Software Foundations]

Automatic Generalisation of Induction Hypotheses

$$\forall n\ m, n + n = m + m \rightarrow n = m$$

...

$$\text{ih} : n + n = m + m \rightarrow n = m$$

$$h : n + n + 2 = m + m$$

$$\vdash n + 1 = m$$

Automatic Generalisation of Induction Hypotheses

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$$\text{ih} : \forall m, n + n = m + m \rightarrow n = m$$

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Naming

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$\forall \alpha (r : \alpha \rightarrow \alpha \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

Naming

$\forall \alpha (r : \alpha \rightarrow \alpha \rightarrow \text{Type}) (a \ b \ c : \alpha)$
 $(h_1 : \text{tc } r \ a \ b) (h_2 : \text{tc } r \ b \ c), \text{tc } r \ a \ c$

$\alpha : \text{Type}$

$r : \alpha \rightarrow \alpha \rightarrow \text{Type}$

$c \ a \ b \ h_{1_x} \ h_{1_y} \ h_{1_z} : \alpha$

$h_{1_hr} : r \ h_{1_x} \ h_{1_y}$

$h_{1_ht} : \text{tc } r \ h_{1_y} \ h_{1_z}$

$h_{1_ih} : \text{tc } r \ h_{1_z} \ c \rightarrow \text{tc } r \ h_{1_y} \ c$

$h_2 : \text{tc } r \ h_{1_z} \ c$

$\vdash \text{tc } r \ h_{1_x} \ c$

Naming

$\forall \alpha (r : \alpha \rightarrow \alpha \rightarrow \text{Type}) (a\ b\ c : \alpha)$
 $(h_1 : \text{tc } r\ a\ b) (h_2 : \text{tc } r\ b\ c), \text{tc } r\ a\ c$

$\alpha : \text{Type}$
 $r : \alpha \rightarrow \alpha \rightarrow \text{Type}$
 $a\ y\ b\ c : \alpha$
 $hr : r\ a\ y$
 $h_1 : \text{tc } r\ y\ b$
 $ih : \forall c, \text{tc } r\ b\ c \rightarrow \text{tc } r\ y\ c$
 $h_2 : \text{tc } r\ b\ c$
 $\vdash \text{tc } r\ a\ c$

Summary

- ▶ Simplification of induction hypotheses
- ▶ Automatic generalisation of induction hypotheses
- ▶ Naming
- ▶ (Not in this talk:) Evaluation of Lean 3's metaprogramming framework

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